

Journées MAS 2022, Rouen

An alternative pricing method for financial models with transaction costs

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Computable prices in model with transaction costs

We consider a portfolio process : $(V_t)_{t=-1}^T$ where $V_{-1} \in \mathbf{R}e_1$ is the initial endowment expressed in cash.

The **solvency set** $\mathbf{G}_t : \Omega \rightarrow \mathbf{R}^d$ is \mathcal{F}_t -measurable random closed set, i.e.

$$\text{Graph}(\mathbf{G}_t) = \{(\omega, x) : x \in \mathbf{G}_t(\omega)\} \in \mathcal{F}_t \otimes \mathcal{B}(\mathbf{R}^d)$$

The **cost process** $\mathbf{C} = (\mathbf{C}_t)_{t=0}^T$ associated to \mathbf{G} is defined as :

$$\mathbf{C}_t(z) = \inf\{\alpha \in \mathbf{R} : \alpha e_1 - z \in \mathbf{G}_t\} = \min\{\alpha \in \mathbf{R} : \alpha e_1 - z \in \mathbf{G}_t\}$$

\mathbf{C}_t is lower semicontinuous and is allowed to be **non-convex**.

Similarly, we may define the liquidation value process $L = (L_t)_{t=0}^T$:

$$L_t(z) := \sup\{\alpha \in \mathbf{R} : z - \alpha e_1 \in \mathbf{G}_t\}, \quad z \in \mathbf{R}^d.$$

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A **portfolio process** is a stochastic process $(V_t)_{t=-1}^T$ satisfying :

$$\Delta V_t = V_t - V_{t-1} \in -\mathbf{G}_t, \text{ a.s.}, \quad t = 0, \dots, T.$$

or equivalently, $L_t(-\Delta V_t) \geq 0$ a.s.

Some examples :

$$\mathbf{C}_t(z) = zS_t, \text{ no transaction cost}$$

$$\mathbf{C}_t(z) = z^1 + S_t^a z^2 \mathbf{1}_{z^2 > 0} + S_t^b z^2 \mathbf{1}_{z^2 < 0}, \text{ bid - ask}$$

and the model with fixed cost (non-convex cost) :

$$\mathbf{C}_t(z) = z^1 - (-z^2 S_t^b - c_t)^+ \mathbf{1}_{z^2 < 0} + (z^2 S_t^a - c_t) \mathbf{1}_{z^2 > 0}$$

We denote by $\mathcal{R}_t(\xi)$ the set of all portfolio processes starting at time $t \leq T$ that replicates ξ at the terminal date T :

$$\mathcal{R}_t(\xi) := \{(V_s)_{s=t}^T, -\Delta V_s \in L^0(\mathbf{G}_s, \mathcal{F}_s), \forall s \geq t+1, V_T = \xi\}.$$

The set of replicating prices of ξ at time t is

$$\mathcal{P}_t(\xi) := \left\{ V_t = (V_t^1, V_t^{(2)}) : (V_s)_{s=t}^T \in \mathcal{R}_t(\xi) \right\}.$$

The infimum replicating cost is then defined as :

$$c_t(\xi) := \text{ess inf}_{\mathcal{F}_t} \{ \mathbf{C}_t(V_t), V_t \in \mathcal{P}_t(\xi) \}.$$

We are interested in $c_0(\xi)$

Definition (Conditional essential infimum)

Let \mathcal{H} and \mathcal{F} be complet σ -algebras such that $\mathcal{H} \subseteq \mathcal{F}$ and let $\Gamma = (\gamma_i)_{i \in I}$ be a family of real-valued \mathcal{F} -measurable random variables. There exists a unique (up to a P -negligible set) random variables $\gamma_{\mathcal{H}} \in L^0(\overline{\mathbf{R}}, \mathcal{H})$, denoted by $\text{ess inf}_{\mathcal{H}} \Gamma$, which satisfies the following properties

- 1) For every $i \in I$, $\gamma_{\mathcal{H}} \leq \gamma_i$ a.s.
- 2) If $\zeta \in L^0(\overline{\mathbf{R}}, \mathcal{H})$ satisfies $\zeta \leq \gamma_i$ a.s. for all $i \in I$, then $\zeta \leq \gamma_{\mathcal{H}}$ a.s.

In classical approach, using dual characterization of $c_0(\xi)$ via [Consistent Price Systems](#).

Numerical algorithm is available only for [finite](#) probability space.

Important observation : $L_T(-\Delta V_T) \geq 0$ is equivalent to :

$$\begin{aligned} V_{T-1}^1 &\geq \xi^1 + \mathbf{C}_T((0, \xi^{(2)} - V_{T-1}^{(2)})), \\ \iff V_{T-1}^1 &\geq \text{ess sup}_{\mathcal{F}_{T-1}} \left(\xi^1 + \mathbf{C}_T((0, \xi^{(2)} - V_{T-1}^{(2)})) \right), \text{ a.s.} \end{aligned}$$

We continue,

$$V_t^1 \geq \text{ess sup}_{\mathcal{F}_t} \left(\xi^1 + \sum_{s=t+1}^T \mathbf{C}_s(0, V_s^{(2)} - V_{s-1}^{(2)}) \right).$$

Computable prices in model with transaction costs

We set :

$$\Pi_u^T(\xi_{u-1}, \xi) := \{\xi_{u-1}^{(2)}\} \times \Pi_{s=u}^{T-1} L^0(\mathbf{R}^{d-1}, \mathcal{F}_s) \times \{\xi^{(2)}\}$$

We define $\gamma_t^\xi(V_{t-1})$ as :

$$\gamma_t^\xi(V_{t-1}) := \operatorname{ess\,inf}_{V^{(2)} \in \Pi_t^T(V_{t-1}, \xi)} \operatorname{ess\,sup}_{\mathcal{F}_t} \left(\xi^1 + \sum_{s=t}^T \mathbf{C}_s(0, V_s^{(2)} - V_{s-1}^{(2)}) \right).$$

Theorem (Dynamic Programming Principle)

For any $0 \leq t \leq T - 1$ and $V_{t-1} \in L^0(\mathbf{R}^d, \mathcal{F}_{t-1})$, we have

$$\gamma_t^\xi(V_{t-1}) = \operatorname{ess\,inf}_{V_t \in L^0(\mathbf{R}^d, \mathcal{F}_t)} \operatorname{ess\,sup}_{\mathcal{F}_t} \left(\mathbf{C}_t(0, V_t^{(2)} - V_{t-1}^{(2)}) + \gamma_{t+1}^\xi(V_t) \right).$$

In particular, $\gamma_t^\xi(0) = c_0(\xi)$.

Usually, we have $\gamma_t^\xi(V_{t-1}) = \gamma_t^\xi(S_t, V_{t-1})$. We consider a weak NA condition :

Definition (Absence of Immediate Profit : AIP)

For any $t \leq T - 1$, the minimal cost to hedge zero payoff is identical to 0 a.s. : $c_t(0) = 0$.

Definition (Strong Absence of Immediate Profit : SAIP)

For any $t \leq T - 1$, the minimal cost to hedge zero payoff is identical to 0 a.s. : $c_t(0) = 0$ (AIP holds). Moreover, $c_t(0)$ is attained only at the zero strategy.

Theorem (ω -wise formulation of DPP)

Under SAIP, the DPP is computable ω -wise as :

$$\gamma_t^\xi(S_t, V_{t-1}) = \inf_{y \in \mathbf{R}^d} \left(\mathbf{C}_t(S_t, (0, y^{(2)} - V_{t-1}^{(2)})) + \sup_{s \in \phi_t(S_t)} \gamma_{t+1}^\xi(s, y) \right),$$

where $\phi_t(S_t) = \text{supp}_{\mathcal{F}_t} S_{t+1}$. Also, the infimum hedging cost of ξ at any time t is reached.

Main idea : maintain the lower semicontinuity of γ_t^ξ .

Theorem (Reachability set)

Under SAIP, the DPP is computable ω -wise as :

$$\gamma_t^\xi(S_t, V_{t-1}) = \inf_{y \in K_t(S_t, V_{t-1})} \left(\mathbf{C}_t(S_t, (0, y^{(2)} - V_{t-1}^{(2)})) \right. \\ \left. + \sup_{s \in \phi_t(S_t)} \gamma_{t+1}^\xi(s, y) \right),$$

for some compact set-valued mapping $K_t : \mathbf{R}^d \times \mathbf{R}^d \rightarrow \mathbf{R}^d$.

We consider Binomial model $\text{supp}_{\mathcal{F}_t} S_{t+1} = \{k_t^d S_t, k_t^u S_t\}$, where $k_t^d, k_t^u \in \mathbf{R}_+$ and the cost function :

$$\mathbf{C}_t(S_t, v) = v^1 + S_t^a v^2 \mathbf{1}_{v^2 \geq 0} + S_t^b v^2 \mathbf{1}_{v^2 \leq 0}.$$

We want to hedge the payoff $\xi = ((K - S_T)^+, 0) \in \mathbf{R}^2$.

We suppose that $S_t^a = S_t(1 + \epsilon_t)$, $S_t^b = S_t(1 - \epsilon_t)$, $\epsilon_t \in \mathbf{R}_+$.

Theorem

AIP_{T-1} holds if and only if :

$$k_{T-1}^d \leq \frac{1 + \epsilon_{T-1}}{1 - \epsilon_T} \text{ and } k_{T-1}^u \geq \frac{1 - \epsilon_{T-1}}{1 + \epsilon_T}.$$

Moreover, $SAIP_{T-1}$ holds if and only if the above inequalities are strict. Moreover, suppose that $1 + \epsilon_{T-1} \leq (1 + \epsilon_T)k_{T-1}^u$ and $1 - \epsilon_{T-1} \geq (1 - \epsilon_T)k_{T-1}^d$, AIP_{T-2} holds if and only if :

$$k_{T-2}^d \leq \frac{1 + \epsilon_{T-2}}{1 + \epsilon_{T-1}} \text{ and } k_{T-2}^u \geq \frac{1 - \epsilon_{T-2}}{(1 - \epsilon_T)k_{T-1}^d}.$$

Moreover, $SAIP_{T-2}$ holds if and only if the above inequalities are strict

Bid-Ask spread

We suppose $T = 2$, the proportional cost coefficients $\epsilon_1 = \epsilon_2 = 0.02$. We assume that SAIP condition holds and choose $k_2^d = 0.9, k_2^u = 1.1, k_1^d = 0.9, k_1^u = 1.2$.

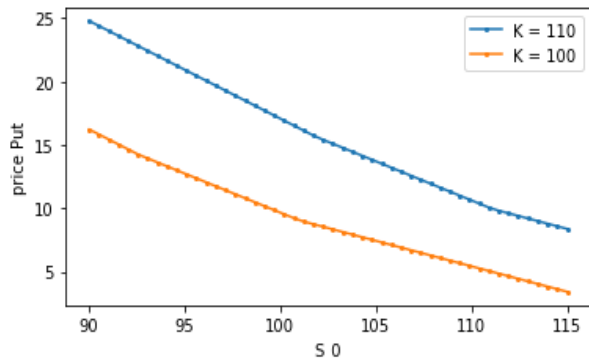


Figure – Price of Put option

Consider the cost function :

$$\mathbf{C}_t(z) = z^1 - (-z^2 S_t^b - c_t)^+ \mathbf{1}_{z^2 < 0} + (z^2 S_t^a - c_t) \mathbf{1}_{z^2 > 0}$$

We suppose that the market defined by \mathbf{C}_t with $c_t = 0$ satisfies SAIP. We called this condition Robust SAIP or RSAIP.

Fixed cost

We use the same parameters as Bid-Ask spread and we consider fixed costs $c_1 = c_2 = 0.8$.

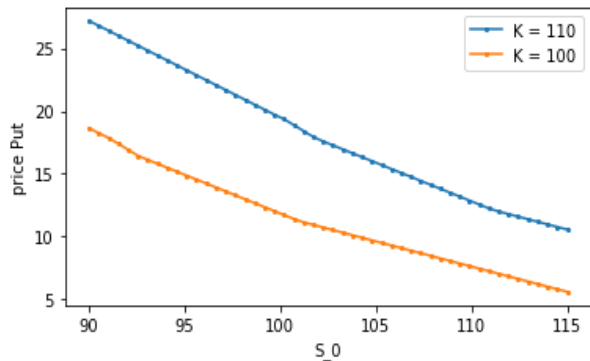




Figure – Price of put option with fixed costs.

-  Lépinette, Emmanuel and Duc-Thinh, Vu. Dynamic programming principle and computable prices in financial market models with transaction costs. Preprint, 2021.
-  Lépinette, Emmanuel and Duc-Thinh, Vu. Limit theorems for the super-hedging prices in general models with transaction costs. Preprint, 2022.

Thank you for your attention !

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